

# HIGHER-ORDER ALGEBRAIC THEORIES AND RELATIVE MONADS

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(Masaryk University Algebra Seminar, 13.05.21)

# Outline

- Algebraic theories
- Second-order algebraic theories
- Higher-order algebraic theories
- A universal characterisation of Law
- Relative monads, monads, and theories
- $0^{\text{th}}$ -order algebraic theories

# I. ALGEBRAIC THEORIES

# First-order operators

$$1. \quad \frac{\Gamma \vdash a \quad \Gamma \vdash b}{\Gamma \vdash a \times b}$$

Multiplication

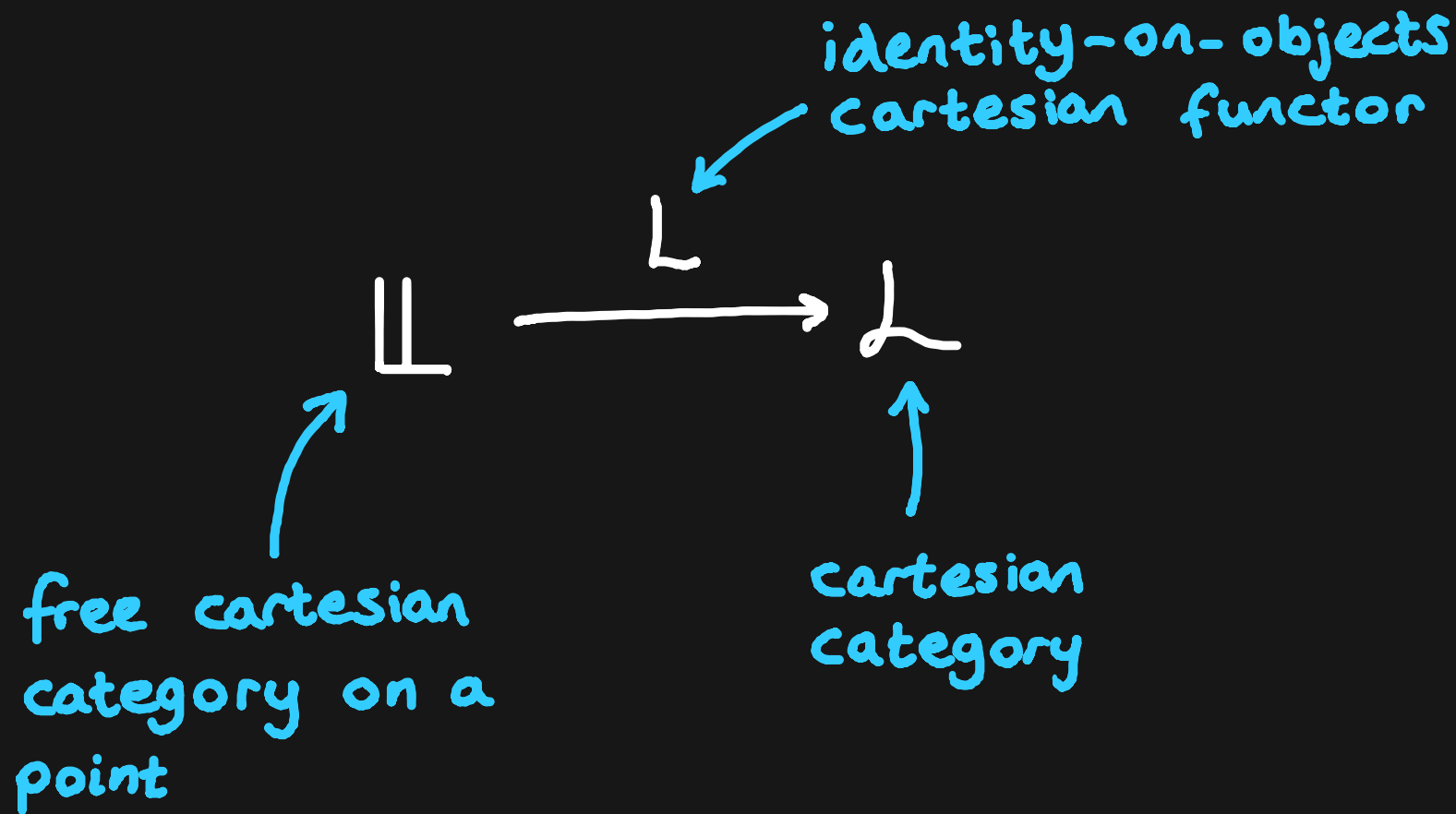
$$2. \quad \frac{\Gamma \vdash a}{\Gamma \vdash a^{-1}}$$

Inverses

$$3. \quad \frac{\Gamma \vdash m : M \quad \Gamma \vdash a : A}{\Gamma \vdash m * a : A}$$

Actions

# Algebraic theories



(Here, 'cartesian' means finite products.)

# Algebraic theories

$$\mathbb{L} \xrightarrow{\quad \mathcal{L} \quad} \mathcal{L}$$

The objects of  $\mathcal{L}$  are given by  $X^n$  for  $X$  the generating object, and  $n \in \mathbb{N}$ .

A morphism  $X^n \xrightarrow{\vec{t}} X^m$  represents an  $m$ -tuple of terms in  $n$  variables:

$$\langle x_1, \dots, x_n + t_i \rangle_{1 \leq i \leq m}$$

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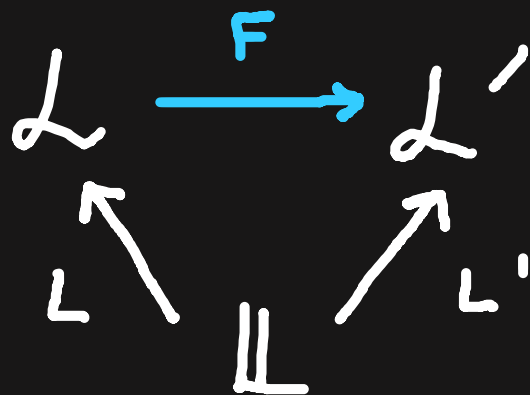
A morphism  $X^n \xrightarrow{\vec{t}} X^m$  represents an  $m$ -tuple of terms in  $n$  variables:

$$\langle x_1, \dots, x_n \vdash t_i \rangle_{1 \leq i \leq m}$$

$n$ -ary operation  $\langle t_i : X^n \rightarrow X \rangle_i$

## Algebraic theories

A **map** of algebraic theories is a commutative triangle



Algebraic theories and their maps form a category

**Law.**



# Law

Law is a well-behaved category.

- It has all small limits and colimits.
- More than that, it is locally strongly finitely presentable:

# Law

Law is a well-behaved category.

- It has all small limits and colimits.
- More than that, it is locally strongly finitely presentable:

$$\text{Law} \simeq \text{Sind}(\mathcal{C}) \simeq \text{Cart}(\mathcal{C}^{\circ}, \text{Set})$$

for some cocartesian  $\mathcal{C}$

Sind = cocompletion under sifted colimits

## Monads and theories

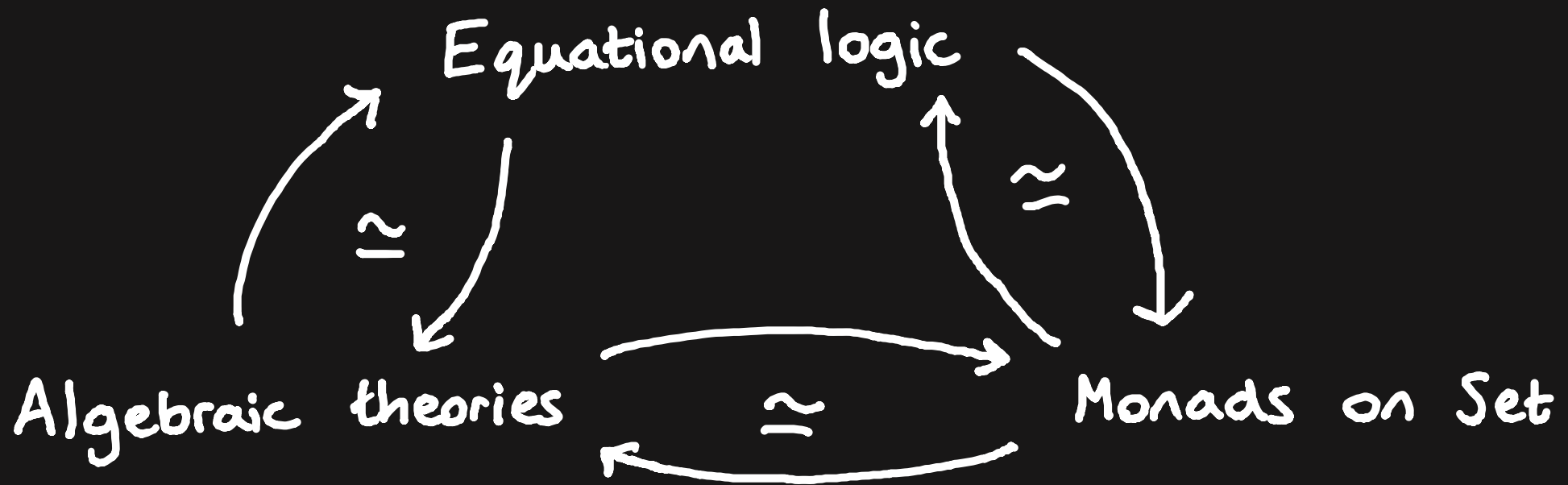
There is a classic equivalence between algebraic theories and (strongly) finitary monads on the category of sets.

$$\text{Law} \simeq \text{Mnd}_f(\text{Set}) = \text{Mnd}_{sf}(\text{Set})$$

Finitary = preserves filtered colimits

Strongly finitary = preserves sifted colimits  
(sifted-cocontinuous)

# Universal algebra



## II. SECOND-ORDER ALGEBRAIC THEORIES

[Fiore & Mahmoud, 2010]

## Second-order operators

1. 
$$\frac{\Gamma, x \vdash f \quad \Gamma \vdash x_0}{\Gamma \vdash \frac{df}{dx}(x_0)}$$
 Differential operators  
(cf. Plotkin 2020)

2. 
$$\frac{\Gamma, x \vdash P}{\Gamma \vdash \exists x. P}$$
 Logical quantifiers

3. 
$$\frac{\Gamma, x \vdash t}{\Gamma \vdash \lambda x. t}$$
  $\lambda$ -abstraction

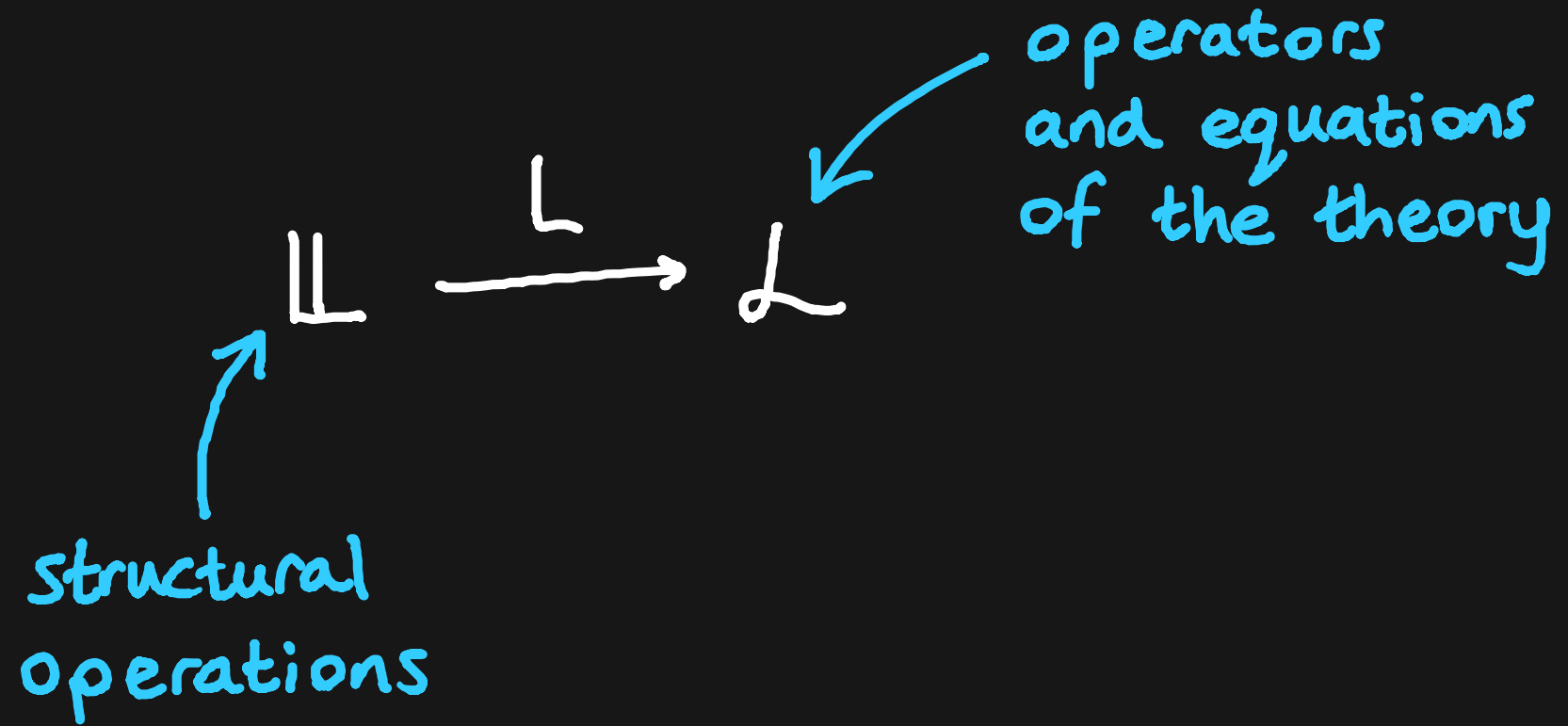
## Second-order operators

$$4. \quad \frac{\Gamma \vdash t : A + B \quad \Gamma, a : A \vdash u : C \quad \Gamma, b : B \vdash v : C}{\Gamma \vdash \text{case}(t, a.u, b.v) : C}$$

Coproducts,  
case-splitting

$$5. \quad \frac{\Gamma, x : X \vdash f : X}{\Gamma \vdash \text{fix}(f) : X} \quad \text{Fixed points}$$

6. Parameterised algebraic theories [Staton, 2013]





## Second-order theory of equality

$\mathbb{L}_2$  is the free cartesian category with an exponentiable object (i.e. an object such that  $(-)^X : \mathbb{L}_2 \rightarrow \mathbb{L}_2$  exists).

Objects of  $\mathbb{L}_2$  are given by products

$$X^{X^{n_1}} \times \dots \times X^{X^{n_k}}$$

with morphisms given by projection and evaluation.

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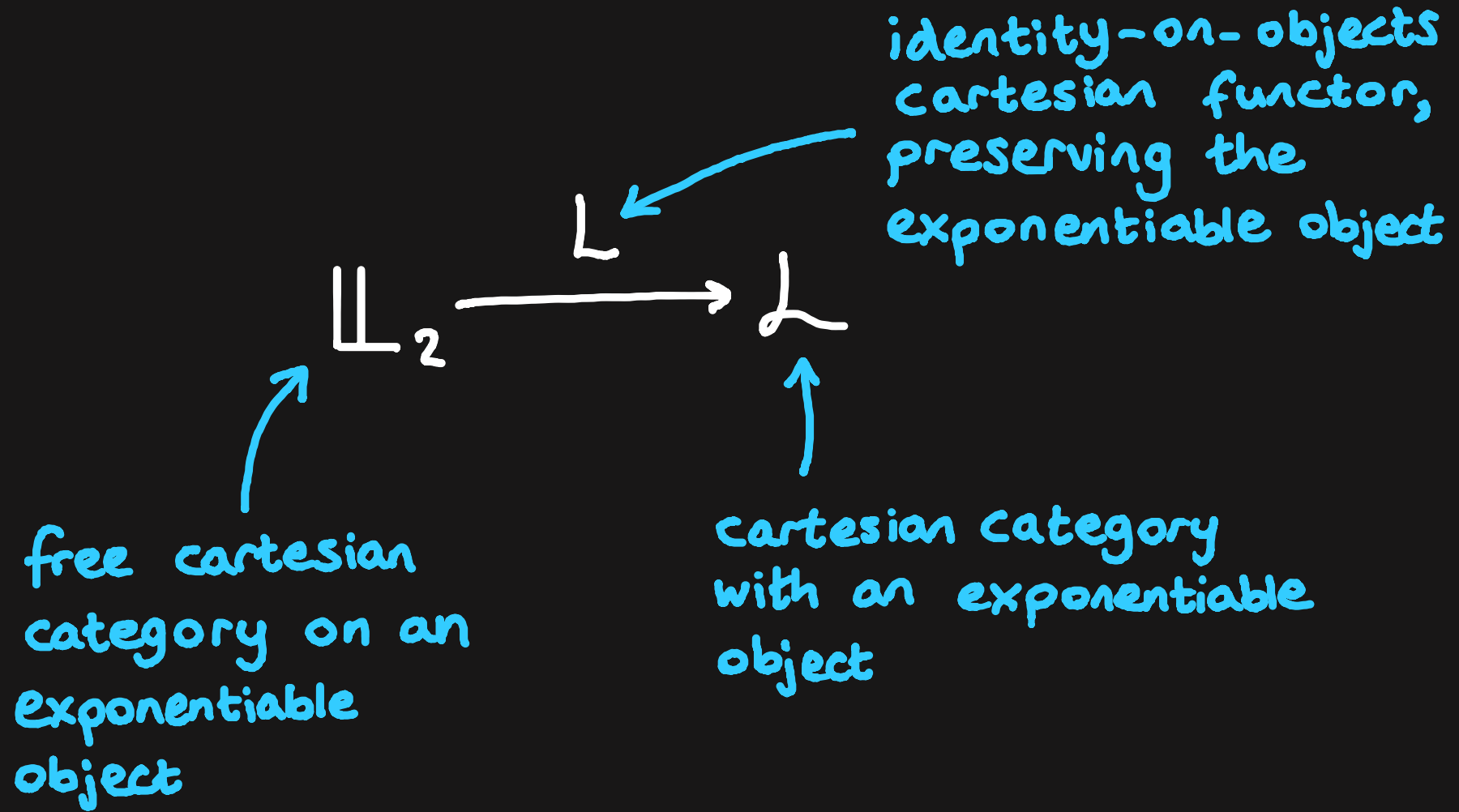
Objects of  $\mathbb{L}_2$  are given by products

$$X^{X^{n_1}} \times \dots \times X^{X^{n_k}}$$

← exponents are  
the objects of  
 $\mathbb{L}$

with morphisms given by projection and evaluation.

# Second-order algebraic theories



## Second-order algebraic theories

$$\mathbb{L}_2 \xrightarrow{\mathcal{L}} \mathcal{L}$$

A morphism  $X^{x^{n_1}} \times \dots \times X^{x^{n_k}} \xrightarrow{t} X^{x^{m_1}} \times \dots \times X^{x^{m_l}}$   
in  $\mathcal{L}$  represents an  $\mathcal{L}$ -tuple of terms in  
 $k$  metavariables and  $m_i$  variables:

$$\langle (x_1^1, \dots, x_{n_1}^1) x_1, \dots, (x_1^k, \dots, x_{n_k}^k) x_k, y_1, \dots, y_{m_i} \vdash t \rangle_i$$

parameterised variable      ordinary variable

'Differentiate  $f(x)$  with respect to  $x$  and evaluate at  $x_0$ .'

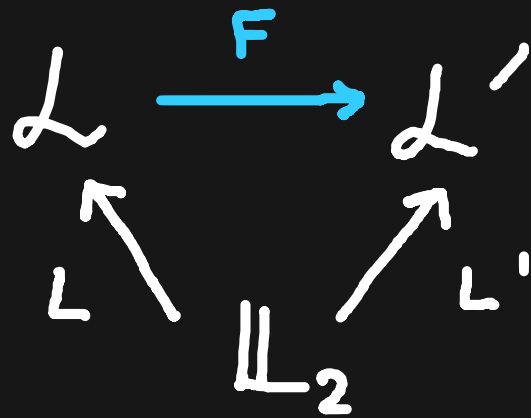
$$\partial(x, f(x), x_0)$$

represented by

$$x^x \times x \xrightarrow{\partial} x$$

## Second-order algebraic theories

A **map** of second-order algebraic theories is a commutative triangle



Second-order algebraic theories and their maps form a category  $\mathbf{Law}_2$ .

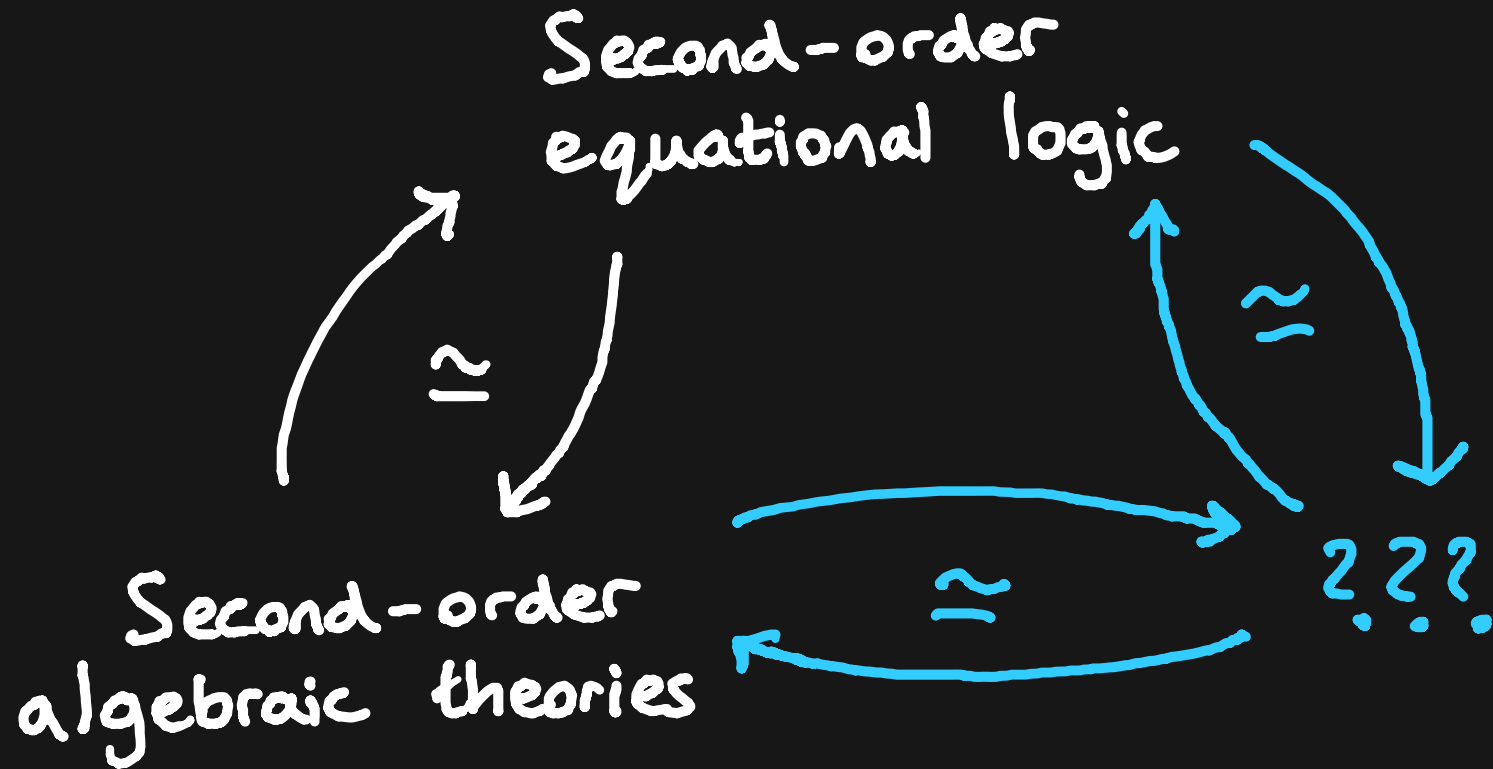
## Second-order algebraic theories

How well-behaved is  $\text{Law}_2$ ?

Does it have:

- Limits?
- Colimits?
- A universal property?

# Second-order universal algebra

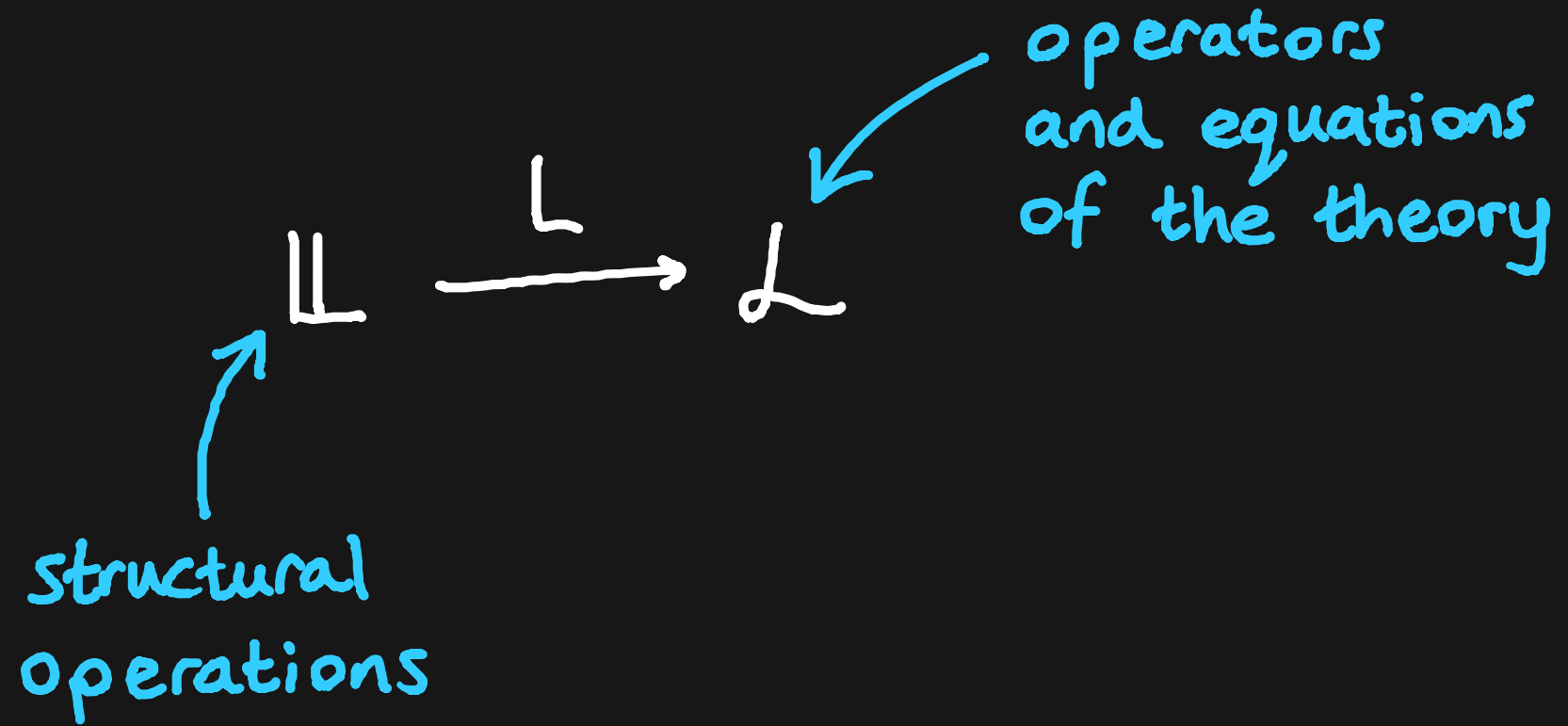




### III. HIGHER-ORDER ALGEBRAIC THEORIES

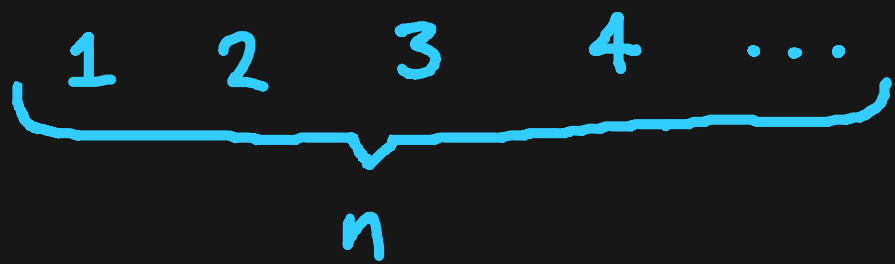
## Third-order operators

1. Continuations
2. Selection operators



## Higher-order theory of equality

$\mathbb{L}_n$  is the free cartesian category on an  $n$ -tetrable object (i.e. an object  $X$  such that  $1, X, X^X, X^{X^X}, \dots$  is exponentiable).



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$\mathbb{L}_n$  is the free cartesian category on an  $n$ -tetrable object (i.e. an object  $X$  such that  $1, X, X^X, X^{X^X}, \dots$  is exponentiable).

$\underbrace{1 \quad 2 \quad 3 \quad 4 \quad \dots}_{n}$

Intuitively, morphisms in  $\mathbb{L}_n$  represent operators taking operators as operands.

## Higher-order theory of equality

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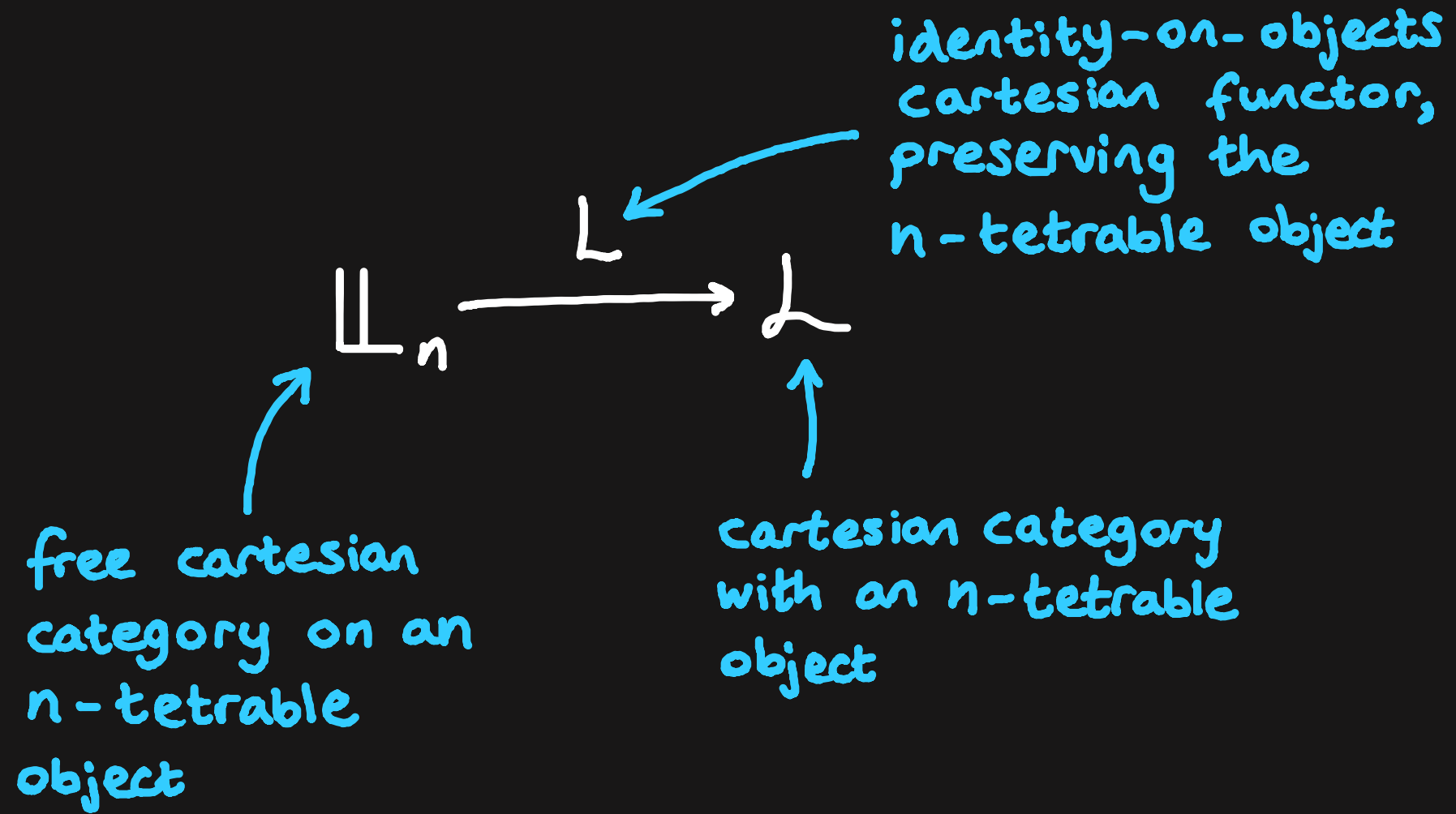
We have:

$$\mathbb{L} = \mathbb{L}_1 \longleftrightarrow \mathbb{L}_2 \longleftrightarrow \dots \longleftrightarrow \mathbb{L}_\omega$$

↑ free cartesian category on a point

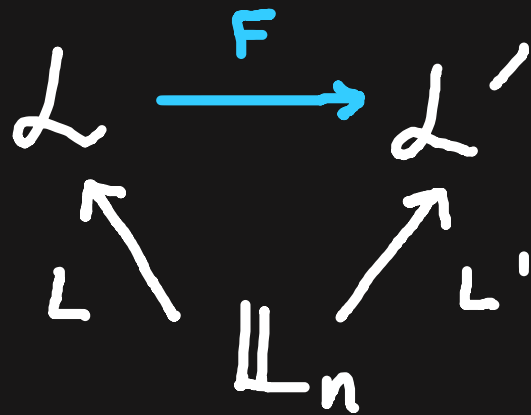
↑ free cartesian-closed category on a point

# Higher-order algebraic theories



## Higher-order algebraic theories

A **map** of  $n^{\text{th}}$ -order algebraic theories is a commutative triangle



$n^{\text{th}}$ -order algebraic theories and their maps form a category  $\mathbb{L}aw_n$ .



1. How well-behaved is  $\text{Law}_n$ ?

2. Is there a monad correspondence?

## Exponentiable subcategories

A (full) subcategory  $\mathcal{C}' \xrightarrow{j} \mathcal{C}$  is exponentiable if, for all  $X \in \mathcal{C}'$ ,  $jX$  is exponentiable in  $\mathcal{C}$ , i.e.  $jX \times (-) \dashv (-)^{jX}$ .

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Ex

$\mathbb{U}_n \hookrightarrow \mathbb{U}_{n+1}$  is exponentiable.

# Exponentiable subcategories

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Denote by  $\text{Exp}_n$  the category with

$$\begin{array}{ccc} \mathcal{C}_1 & \longleftrightarrow & \dots & \longleftrightarrow & \mathcal{C}_n \\ F_1 \downarrow & & \dots & & \downarrow F_n \\ \mathcal{D}_1 & \longleftrightarrow & \dots & \longleftrightarrow & \mathcal{D}_n \end{array}$$

where  $F_{i+1}$  preserves exponentiation by objects of  $\mathcal{C}_i$

Prop.

There is an adjunction

$$\text{Exp}_n \begin{array}{c} \xrightarrow{\Delta^{-1}} \\ \xleftarrow{\Delta} \end{array} \text{Exp}_{n+1}$$

conservative  
cartesian  
closure of  $\mathcal{L}_n$

where

$$\nabla \mathcal{L}_1 \hookrightarrow \dots \hookrightarrow \mathcal{L}_n \nabla = \mathcal{L}_1 \hookrightarrow \dots \hookrightarrow \mathcal{L}_n \dashrightarrow \tilde{\mathcal{L}} \quad \swarrow$$

$$\Delta \mathcal{L}_1 \hookrightarrow \dots \hookrightarrow \mathcal{L}_{n+1} \Delta = \mathcal{L}_1 \hookrightarrow \dots \hookrightarrow \mathcal{L}_n$$

Prop.

Let  $\mathcal{C}_1 \xrightarrow{j} \mathcal{C}_2$  be an object of  $\text{Exp}_2$ .

Strict cartesian functors  $\tilde{\mathcal{C}}_2 \rightarrow \text{Set}$  are in bijection with strict cartesian identity-on-objects functors out of  $\mathcal{C}_2$  preserving exponentiation by objects of  $\mathcal{C}_1$ .

$$\text{Cart}(\tilde{\mathcal{C}}_2, \text{Set}) \cong j /_{i_0} \text{Exp}_2$$

## Prop.

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$$\text{Cart}(\tilde{\mathcal{C}}_2, \text{Set}) \cong j /_{i_0} \text{Exp}_2$$

## Sketch

For  $F: \tilde{\mathcal{C}}_2 \rightarrow \text{Set}$ ,  $F(Y^X)$  defines a hom-set, with composition induced by evaluation.

Cor

$\mathbb{L}_n \xrightarrow{j} \mathbb{L}_{n+1}$  is an object of  $\text{Exp}_2$ , so

strict cartesian functors  $\widetilde{\mathbb{L}}_{n+1} \rightarrow \text{Set}$  are in bijection with strict cartesian identity-on-objects functors out of  $\mathbb{L}_{n+1}$  preserving exponentiation by objects of  $\mathbb{L}_n$ .

$$\text{Cart}(\widetilde{\mathbb{L}}_{n+1}, \text{Set}) \cong j /_{i_0} \text{Exp}_2 = \text{Law}_{n+1}$$

$\uparrow$   
 $\mathbb{L}_{n+2}$

$$\begin{array}{ccc} \mathbb{L}_{n+1} & \longrightarrow & \mathcal{L}' & (A) \\ j \uparrow & & \uparrow & \\ \mathbb{L}_n & \longrightarrow & \mathcal{L} & (B) \end{array}$$



# The universal property of $\mathcal{L}aw_n$

Thm

$\mathcal{L}aw_n$  is locally strongly finitely presentable.

$$\mathcal{L}aw_n \simeq \text{Cart}(\mathcal{U}_{n+1}, \text{Set})$$

$$\text{sifted cocompletion} \simeq \text{Sind}(\mathcal{U}_{n+1}^\circ)$$

free cartesian category on an  $(n+1)$ -tetrable point

# The universal property of Law<sub>n</sub> (n=1)

Thm

Law<sub>1</sub> is locally strongly finitely presentable.

$$\text{Law} = \text{Law}_1 \simeq \text{Cart}(\mathbb{U}_2, \text{Set})$$

$$\text{sifted cocompletion} \simeq \text{Sind}(\mathbb{U}_2^{\circ})$$

free cartesian category on an exponentiable object

# The universal property of $\mathcal{L}aw_n$

Thm

$\mathcal{L}aw_n$  is locally strongly finitely presentable.

$$\mathcal{L}aw_n \simeq \text{Cart}(\mathcal{U}_{n+1}, \text{Set})$$

sifted cocompletion  $\curvearrowright$   $\simeq \text{Sind}(\mathcal{U}_{n+1}^\circ)$   $\curvearrowright$  free cartesian category on an  $(n+1)$ -tetrable point

Hence also:

- Locally finitely presentable
- Cocomplete
- Complete

# The universal property of $\text{Law}_n$

Thm

$\text{Law}_n$  is locally strongly finitely presentable.

$$\text{Law}_n \simeq \text{Cart}(\mathbb{L}_{n+1}, \text{Set})$$

sifted cocompletion  $\simeq \text{Sind}(\mathbb{L}_{n+1}^\circ)$  free cartesian category on an  $(n+1)$ -tetrable point

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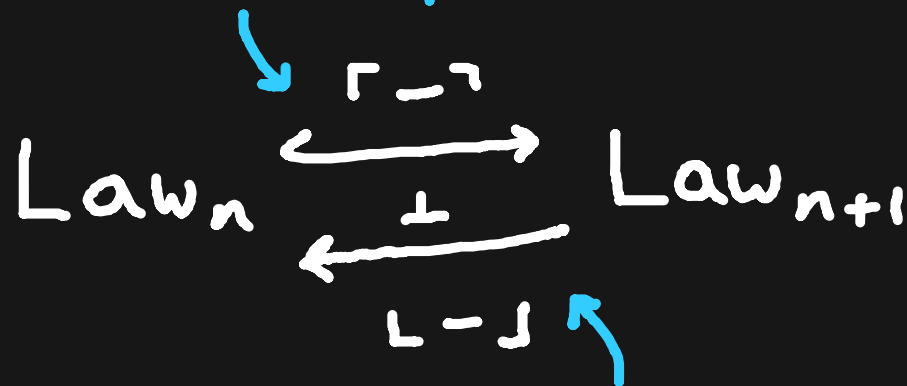
(Cf. Uemura 2020, 'The universal exponentiable arrow'.)

# Coreflections

Thm

There is a coreflection of categories:

inclusion of presentations



discard  $(n+1)^{\text{th}}$ -order terms

# Coreflections

Thm

There is a coreflection of categories

$$\text{Sind}(\mathbb{L}_{n+1}) \cong \text{Law}_n \begin{array}{c} \xrightarrow{\Gamma \dashv} \\ \xleftarrow{\perp} \\ \xleftarrow{\text{L} \dashv} \end{array} \text{Law}_{n+1} \cong \text{Sind}(\mathbb{L}_{n+2})$$

$\mathbb{L}_{n+1} \xrightarrow{j} \mathbb{L}_{n+2}$  induces the algebraic functor

$\text{L} \dashv$ , which has a left adjoint by abstract nonsense.

## Coreflections

There is a chain of coreflections,

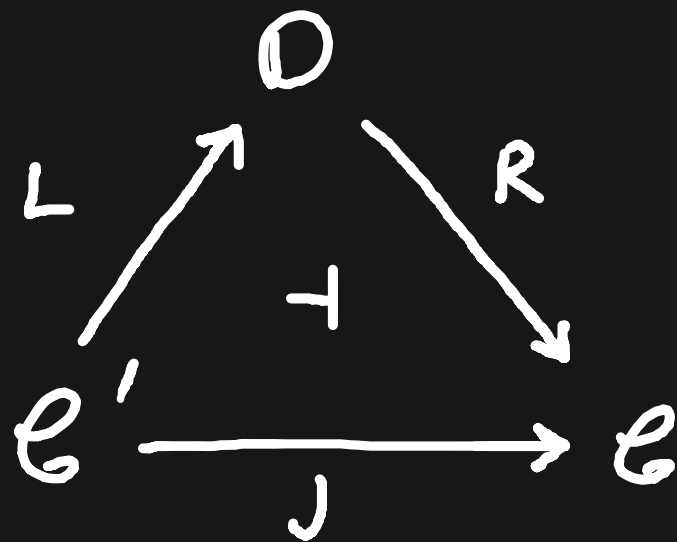
$$\text{Law}_1 \begin{array}{c} \longleftarrow \\ \perp \\ \longrightarrow \end{array} \text{Law}_2 \begin{array}{c} \longleftarrow \\ \perp \\ \longrightarrow \end{array} \dots \begin{array}{c} \longleftarrow \\ \perp \\ \longrightarrow \end{array} \text{Law}_\omega$$

allowing us to freely extend or restrict the order of a higher-order algebraic theory.

## IV. RELATIVE MONADS



## Relative adjunctions



$L \dashv J \vdash R$  when  $\mathcal{D}(Lx, y) \cong \mathcal{C}(Jx, Ry)$  natural  
in  $x \in \mathcal{C}'$  and  $y \in \mathcal{D}$ .

## Relative monads

A  $J$ -relative monad  $(T, \eta, (-)^*)$  consists of

- a function  $T: |E'| \rightarrow |E|$

- a transformation  $\eta_x: JX \rightarrow TX$

- a transformation  $(-)^*_{x,y}: \mathcal{E}(JX, TY) \rightarrow \mathcal{E}(TX, TY)$

satisfying unitality and associativity conditions.

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Prop: a monad is precisely an  $Id$ -relative monad.

Prop: every relative adjunction induces a relative monad.

## Simple slice category

Let  $L: \mathcal{L}_n \rightarrow \mathcal{L}$  be an  $n^{\text{th}}$ -order algebraic theory.

Define a functor

$$L//(-) : \mathcal{L}^{\circ} \longrightarrow \mathcal{L}aw_n$$

where  $L//X$  is the simple slice category over  $X$ :

$$L//X(A, B) = \mathcal{L}(A \times X, B)$$

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where  $L//X$  is the **simple slice category** over  $X$ :

$$L//X(A, B) = \mathcal{L}(A \times X, B)$$

which is the free cartesian category with a morphism  $1 \rightarrow X$  containing  $d$ .

## Theories to relative monads

Lem

Let  $L: \mathbb{L}_{n+1} \rightarrow \mathcal{L}$  be an  $(n+1)^{\text{th}}$ -order algebraic theory.  
There is a  $\mathcal{L}$ -relative adjunction:

$$\begin{array}{ccc} & \mathcal{L}^{\circ} & \\ L^{\circ} \nearrow & & \searrow \llbracket \mathcal{L} // - \rrbracket \\ \mathbb{L}_{n+1}^{\circ} & \xleftrightarrow{\mathcal{L}} & \text{Sind}(\mathbb{L}_{n+1}^{\circ}) \end{array}$$

Hence  $L$  induces a  $\mathcal{L}$ -relative monad.

## Relative monads to theories

Lem

Let  $T: \mathbb{L}_{n+1}^\circ \longrightarrow \text{Sind}(\mathbb{L}_{n+1}^\circ)$  be a

$\mathfrak{k}$ -relative monad. There is a  $\mathfrak{k}$ -relative adjunction:

$$\begin{array}{ccc} & \text{KI}(T) & \\ \begin{array}{c} \nearrow k \\ \text{KI}(T) \\ \searrow u \end{array} & & \\ \mathbb{L}_{n+1}^\circ & \xleftarrow{\mathfrak{k}} & \text{Sind}(\mathbb{L}_{n+1}^\circ) \end{array}$$

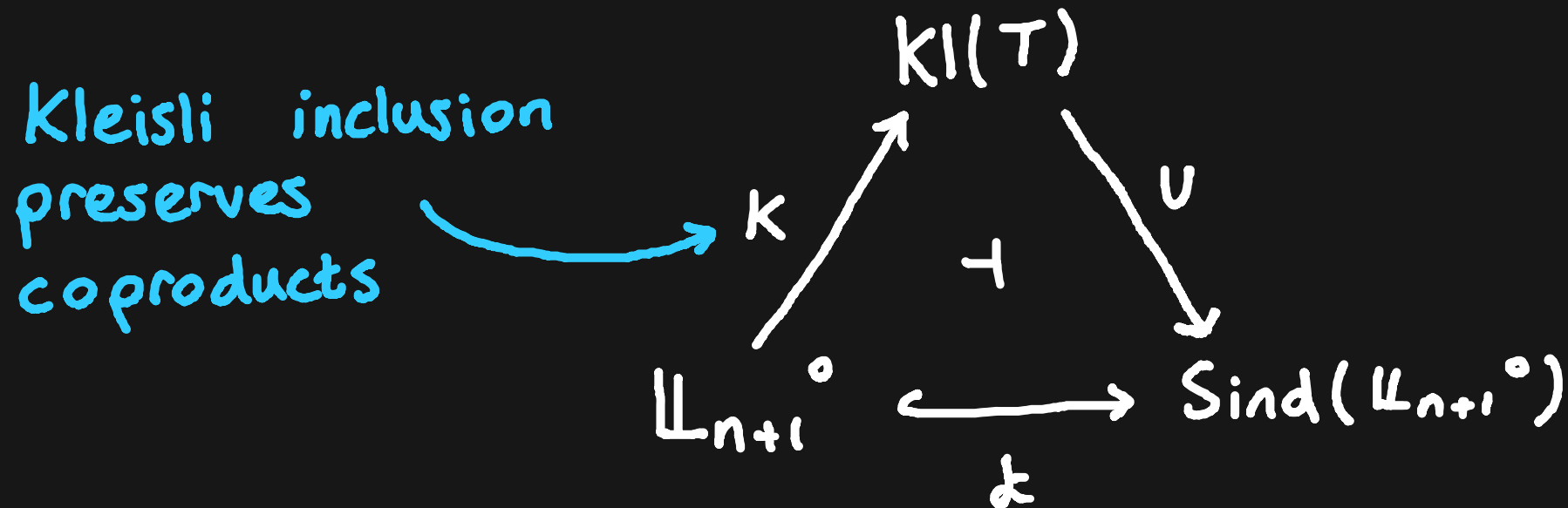


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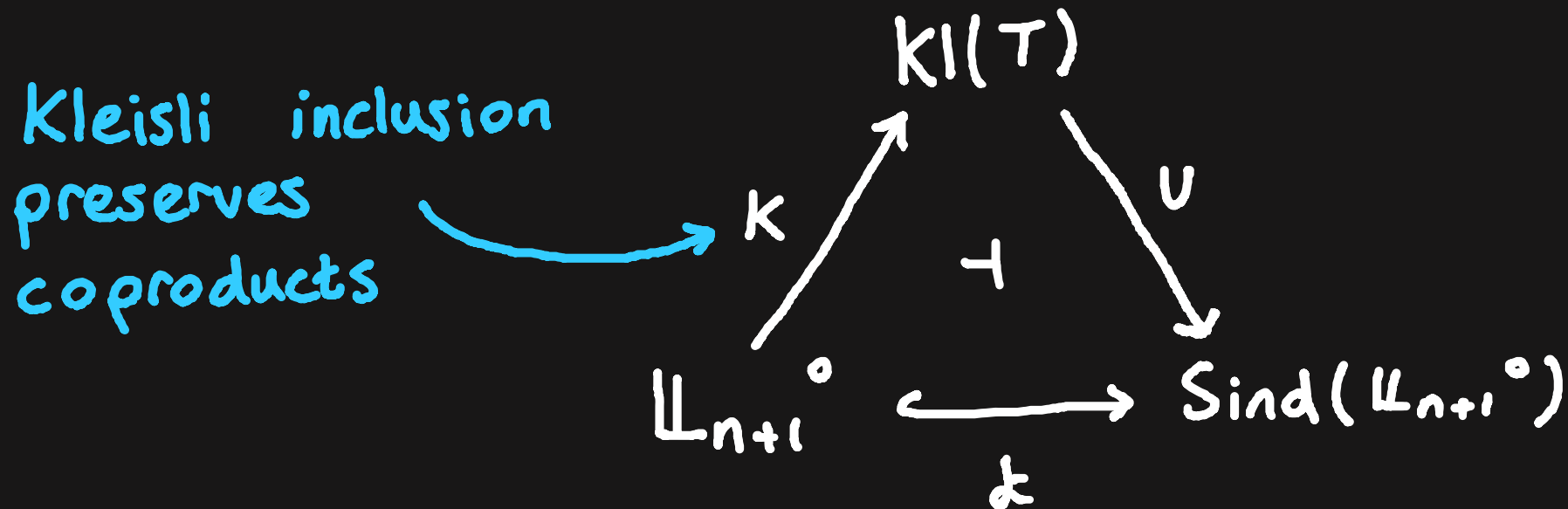


$K^\circ$  is identity-on-objects and cartesian...

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$K^\circ$  is identity-on-objects and cartesian...  
but may not preserve exponentials ☹

## Interlude

When does the Kleisli inclusion preserve coexponentials?

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$$\begin{aligned} & \text{kl}(T)(x, y + z) \\ &= \text{Law}_n(x, T(y + z)) \\ &\stackrel{?}{\cong} \text{Law}_n(x, y + Tz) \\ &\cong \text{Law}_n(x_y, Tz) \\ &= \text{kl}(T)(x_y, z) \end{aligned}$$

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Hence we require

$$T(Y + Z) \cong Y + TZ \quad (Y, Z \in \mathbb{L}_{n-1})$$

Call such relative monads  
**+ -linear**.

## Interlude

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Hence we require

$$T(Y + Z) \cong Y + TZ \quad (Y, Z \in \mathbb{L}_{n-1})$$

Call such relative monads  
**+ - linear.**

(+ - linearity is trivial for  
 $n=1$ .)

## Relative monads to theories

Lem

Let  $T: \mathbb{L}_{n+1}^\circ \longrightarrow \text{Sind}(\mathbb{L}_{n+1}^\circ)$  be a

$\mathfrak{k}$ -relative monad. There is a  $\mathfrak{k}$ -relative adjunction:

$$\begin{array}{ccc} & \mathbb{K}l(T) & \\ \mathfrak{k} \nearrow & & \searrow \mathfrak{U} \\ \mathbb{L}_{n+1}^\circ & \xleftrightarrow{\quad \mathfrak{T} \quad} & \text{Sind}(\mathbb{L}_{n+1}^\circ) \\ & \mathfrak{k} \longleftarrow & \end{array}$$

## Relative monads to theories

Lem

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## Relative monads to theories

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Let  $T: \mathbb{L}_{n+1}^\circ \longrightarrow \text{Sind}(\mathbb{L}_{n+1}^\circ)$  be a  $+ -$ linear

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$K^\circ: \mathbb{L}_{n+1}^\circ \longrightarrow \text{KI}(T)^\circ$  is an  $(n+1)^{\text{th}}$ -order algebraic theory.

Thm

$$\text{Law}_{n+1} \cong \text{RMnd}_{+-\text{lin}}(\mathcal{K}_{\mathbb{L}_{n+1}^\circ})$$

$(n+1)^{\text{th}}$ -order algebraic theories

$+-\text{linear } (\mathbb{L}_{n+1}^\circ \longleftrightarrow \text{Law}_n) -$   
relative monads

Thm

$$\mathbf{Law}_{n+1} \cong \mathbf{RMnd}_{+-\text{lin}}(\mathcal{K}_{\mathbb{L}_{n+1}^\circ})$$

$(n+1)^{\text{th}}$ -order algebraic theories

$+-\text{linear } (\mathbb{L}_{n+1}^\circ \longleftrightarrow \mathbf{Law}_n) \text{-relative monads}$

(But what about ordinary monads?)

## Relative monads & monads

Thm

Let  $\mathcal{C}$  be a locally strongly finitely presentable category. There is an equivalence of categories

$$\text{RMnd}(\mathcal{C}_{\text{sfp}} \hookrightarrow \mathcal{C}) \simeq \text{Mnd}_{\text{sf}}(\mathcal{C})$$

Thm

$$\begin{aligned} \text{Law}_{n+1} &\simeq \text{RMnd}_{+-\text{lin}}(\mathcal{K}\mathbb{U}_{n+1}^\circ) \\ &\text{+-linear } (\mathbb{U}_{n+1}^\circ \longleftrightarrow \text{Law}_n)\text{-} \\ &\text{relative monads} \\ &\simeq \text{Mnd}_{+-\text{lin}, \text{sf}}(\text{Law}_n) \\ &\text{sifted-cocontinuous +-linear} \\ &\text{monads on } \text{Law}_n \end{aligned}$$

$(n+1)^{\text{th}}$ -order algebraic theories

$\text{RMnd}_{+-\text{lin}}(\mathcal{K}\mathbb{U}_{n+1}^\circ)$

$\text{+-linear } (\mathbb{U}_{n+1}^\circ \longleftrightarrow \text{Law}_n)\text{-}$   
 $\text{relative monads}$

$\text{Mnd}_{+-\text{lin}, \text{sf}}(\text{Law}_n)$

$\text{sifted-cocontinuous +-linear}$   
 $\text{monads on } \text{Law}_n$

Prop.  
Let  $L : \mathbb{L}_{n+1} \rightarrow \mathcal{L}$  be an  $(n+1)^{\text{th}}$ -order algebraic theory.  
The corresponding monad is given by

$$T_L(X) = \mathbb{L}L + \lceil X \rceil$$

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The corresponding monad is given by

$$T_L(X) = \mathbb{L}L + \lceil X \rceil$$

When  $n=0$ , this says that  $T_L$  takes a set of constants, freely adds them to  $L$ , then extracts the new constants formed from those in  $X$  under the operations of  $L$ .

Prop.  
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$$T_L(X) = \mathbb{L}L + \ulcorner X \urcorner$$

$$T_L(X)(B, C) \cong \int^{A \in \mathbb{L}_{n+1}} \mathcal{L}(A, C^B) \times \ulcorner X \urcorner(1, A)$$



# Algebras

Let  $L: \mathbb{L}_{n+1} \rightarrow \mathcal{L}$  be an  $(n+1)^{\text{th}}$ -order algebraic theory,  
and let  $T_L: \text{Law}_n \rightarrow \text{Law}_n$  be the corresponding monad.

$$T_L\text{-Alg} \simeq \text{Cart}(\mathcal{L}, \text{Set})$$

## 0<sup>th</sup>-order algebraic theories

It is well-known that (first-order) algebraic theories correspond to (strongly) finitary monads on  $\text{Set}$ .

Since  $\text{Law}_{n+1} \cong \text{Mnd}_{+-\text{lin}, \text{sf}}(\text{Law}_n)$ , and  $+-\text{linearity}$  is trivial when  $n < 2$ , we are led to conclude that

$$\text{Law}_0 \cong \text{Set}$$

How may we interpret this syntactically?

## 0<sup>th</sup>-order algebraic theories

A 0<sup>th</sup>-order algebraic theory is a strict, terminal object-preserving identity-on-objects functor

$$\{1 \leftarrow X\} \longrightarrow \mathbb{L}_0 \xrightarrow{L} \mathcal{L}$$

for which every morphism in  $\mathcal{L}$  is constant (i.e. factors through 1).

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Equivalently, a T-operad for T the terminal monad on Set.

## 0<sup>th</sup>-order algebraic theories

0<sup>th</sup>-order algebraic theories are theories of constants.

Each  $L: \mathbb{L}_0 \rightarrow \mathcal{L}$  defines a set  $\mathcal{L}(1, X)$ , and vice versa, exhibiting an isomorphism of categories

$$\mathbb{L}_0 \cong \text{Set}$$

## 0<sup>th</sup>-order algebraic theories

Recall that:

$$\text{Law}_n \cong \text{Sind}(\mathbb{L}_{n+1}^\circ) \quad (n > 0)$$

Substituting  $n=0$ :

$$\begin{aligned} \text{Law}_0 &\cong \text{Sind}(\mathbb{L},^\circ) \\ &\cong \text{Sind}(\text{FinSet}) \\ &\cong \text{Set} \end{aligned}$$

Everything previously discussed remains valid for  $n=0$ .

Cor.

$$\text{Law}_n \cong \text{Mnd}_{+-\text{lin}, \text{sf}^n}(\text{Law}_0)$$

## Summary

- Higher-order algebraic theories generalise algebraic theories by (higher-order) variable binding operators.
- There are coreflections  $\text{Law}_n \overset{\leftarrow}{\underset{\perp}{\rightleftarrows}} \text{Law}_{n+1}$ .
- $\text{Law}_n \simeq \text{Sind}(\mathbb{L}_{n+1}^\circ)$
- $\text{Law}_{n+1} \simeq \text{Mnd}_{\text{sf}, +\text{-lin}}(\text{Law}_n)$
- Sets are  $0^{\text{th}}$ -order algebraic theories.